

15-213, Fall 2009
Lab Assignment L4: Writing Your Own Unix Shell
Assigned: Tue, Oct 6, Due: Thu, Oct 22, 11:59PM
Last Possible Time to Turn in: Thr, Oct 24, 11:59PM

Ben Blum and Teddy Martin (15-213-staff@cs.cmu.edu) are the lead TAs for this assignment.

1 Introduction

The purpose of this assignment is to become more familiar with the concepts of process control and signalling. You'll do this by writing a simple Unix shell program that supports job control and I/O redirection. Please read the whole writeup before starting.

2 Logistics

This is an individual project. You should do this lab on one of the Fish machines. As always, clarifications and corrections will be posted on the Autolab message board.

3 Hand Out Instructions

Download the file `tshlab-handout.tar` from Autolab, and copy it to the protected directory (the *lab directory*) in which you plan to do your work. Then do the following *on a Fish machine*:

- Type the command `tar xvf tshlab-handout.tar` to expand the tar-file.
- Type your name and Andrew ID in the header comment at the top of `tsh.c`.
- Type the command `make` to compile and link the driver, the trace interpreter, and the test routines.

Looking at the `tsh.c` (*tiny shell*) file, you will see that it contains a skeleton of a simple Unix shell. It will not, of course, function as a shell if you compile and run it now. To help you get started, we have already implemented the less interesting functions such as the routines that manipulate the job list and the command

line parser. Your assignment is to complete the remaining empty functions listed below. As a sanity check for you, we've listed the approximate number of lines of code for each of these functions in our reference solution (which includes lots of comments, this is a good thing).

- `eval`: Main routine that parses and interprets the command line. [300 lines, including some helper functions]
- `sigchld_handler`: Catches SIGCHLD signals. [40 lines]

When you wish to test your shell, type `make` to recompile it. To run it, type `tsh` to the command line:

```
linux> ./tsh
tsh> [type commands to your shell here]
```

4 General Overview of Unix Shells

A **shell** is an interactive command-line interpreter that runs programs on behalf of the user. A shell repeatedly prints a prompt, waits for a **command line** on `stdin`, and then carries out some action, as directed by the contents of the command line.

The command line is a sequence of ASCII text words delimited by whitespace. The first word in the command line is either the name of a built-in command or the pathname of an executable file. The remaining words are command-line arguments:

- If the first word is a built-in command, the shell immediately executes the command in the current process.
- Otherwise, the word is assumed to be the pathname of an executable program. In this case, the shell forks a child process, then loads and runs the program in the context of the child.

The child processes created as a result of interpreting a single command line are known collectively as a **job**. In general, a job can consist of multiple child processes connected by Unix pipes. However, the shell you write in this lab need not support pipes.

If the command line ends with an ampersand “&”, then the job runs in the **background**, which means that the shell does not wait for the job to terminate before printing the prompt and awaiting the next command line. Otherwise, the job runs in the **foreground**, which means that the shell waits for the job to terminate before awaiting the next command line. Thus, at any point in time, at most one job can be running in the foreground. However, an arbitrary number of jobs can run in the background.

For example, typing the command line

```
tsh> jobs
```

causes the shell to execute the built-in `jobs` command. Typing the command line

```
tsh> /bin/ls -l -d
```

runs the `ls` program in the foreground. By convention, the shell ensures that when the program begins executing its main routine

```
int main(int argc, char *argv[])
```

the `argc` and `argv` arguments have the following values:

```
argc == 3
argv[0] == ``/bin/ls``
argv[1]== ``-l``
argv[2]== ``-d``
```

Alternatively, typing the command line

```
tsh> /bin/ls -l -d &
```

runs the `ls` program in the background.

Unix shells support the notion of **job control**, which allows users to move jobs back and forth between background and foreground, and to change the process state (running, stopped, or terminated) of the processes in a job. For example,

- Typing `ctrl-c` causes a `SIGINT` signal to be delivered to each process in the foreground job. The default action for `SIGINT` is to terminate the process.
- Similarly, typing `ctrl-z` causes a `SIGTSTP` signal to be delivered to each process in the foreground job. The default action for `SIGTSTP` is to place a process in the stopped state, where it remains until it is awakened by the receipt of a `SIGCONT` signal.

Unix shells also provide various built-in commands that support job control. For example:

- `jobs`: List the running and stopped background jobs.
- `bg job`: Change a stopped background job into a running background job.
- `fg job`: Change a stopped or running background job into a running foreground job.
- `kill job`: Terminate a job.

Unix shells also support the notion of **I/O redirection**, which allows users to redirect `stdin` and `stdout` to disk files. For example, typing the command line

```
tsh> /bin/ls > foo
```

redirects the output of `ls` to a file called `foo`. Similarly,

```
tsh> /bin/cat < foo
```

displays the contents of file `foo` on `stdout`.

5 The `tsh` Specification

Your `tsh` shell should have the following features:

- The prompt should be the string “`tsh>` ”.
- The command line typed by the user should consist of a name and zero or more arguments, all separated by one or more spaces. If name is a built-in command, then `tsh` should handle it immediately and wait for the next command line. Otherwise, `tsh` should assume that name is the path of an executable file, which it loads and runs in the context of an initial child process (In this context, the term *job* refers to this initial child process). If you are running system programs like `ls`, you will need to enter the full path (in this case `/bin/ls`) because your shell does not have search paths. (Alternatively, if your shell uses `execvp` instead of `execve`, it will search in your `$PATH` for the executable, but this is not required.)

- `tsh` need not support pipes (`|`), but **MUST** support I/O redirection (“`<`” and “`>`”), for instance:

```
tsh> /bin/cat < foo > bar
```

Your shell must support both input and output redirection in the same command line.

- Typing `ctrl-c` should cause a `SIGINT` signal to be sent to the current foreground job (and `ctrl-z` for `SIGTSTP`), as well as any descendants of that job (e.g., any child processes that it forked). If there is no foreground job, then the signal should have no effect.
- If the command line ends with an ampersand `&`, then `tsh` should run the job in the background. Otherwise, it should run the job in the foreground.
- Each job can be identified by either a process ID (PID) or a job ID (JID), which is a positive integer assigned by `tsh`. JIDs should be denoted on the command line by the prefix `%`. For example, “`%5`” denotes JID 5, and “`5`” denotes PID 5. (We have provided you with all of the routines you need for manipulating the job list.)
- `tsh` should support the following built-in commands:
 - The `quit` command terminates the shell.
 - The `jobs` command lists all background jobs.
 - The `bg job` command restarts *job* by sending it a `SIGCONT` signal, and then runs it in the background. The *job* argument can be either a PID or a JID.
 - The `fg job` command restarts *job* by sending it a `SIGCONT` signal, and then runs it in the foreground. The *job* argument can be either a PID or a JID.

- Your shell should be able to redirect the output from the `jobs` built-in command. For example

```
tsh> jobs > foo
```

- `tsh` should reap all of its zombie children. If any job terminates because it receives a signal that it didn’t catch, then `tsh` should recognize this event and print a message with the job’s PID and a description of the offending signal.

6 Process groups and TTY ownership

When you run your shell from the standard Unix shell, your shell is running in the foreground process group. If your shell then creates a child process, by default that child will also be a member of the foreground process group. Since typing `ctrl-c` sends a `SIGINT` to every process in the foreground group, typing `ctrl-c` will send a `SIGINT` to your shell, as well as to every process that your shell created (including background processes!), which obviously isn't correct.

Use of `setpgid`. After the `fork`, but before the `exec`, the child process should call `setpgid(0, 0)`, which puts the child in a new process group whose group ID is identical to the child's PID. Putting your children in different process groups will limit which of them receive signals, and also which of them can read input from the terminal.

Use of `tcsetpgrp`. When you change a *foreground* child process's process group ID with `setpgid`, you will need to also give it ownership of the controlling terminal with `tcsetpgrp`. (Though it's not necessary, consider using `isatty` to find an appropriate file descriptor for `tcsetpgrp`'s first argument. The manpages for `tcsetpgrp`, `isatty`, and `stdin` should prove useful.) When a signal is sent from the terminal, only processes in the process group that owns the terminal will receive that signal. Also, only processes in that process group are allowed to read input from the terminal - any other process attempting to do so will receive `SIGTTIN` and be stopped as though by `SIGSTOP`.

Make sure you understand how use of `setpgid` and `tcsetpgrp` affects which processes have read access to the terminal and which processes receive user-sent signals, and remember to consult the reference shell if there's any confusion about how your shell should behave.

Note. When autolab tests your shell, it will be running in an environment without a controlling tty. This means that your shell will need to behave appropriately even when there is no file descriptor with which you can successfully call `tcsetpgrp`. It also means that autolab can't test for proper usage of `tcsetpgrp` - the TAs will check for that when going over your code for style. Fortunately, it's easy to verify by hand: Make sure that, after implementing use of these two functions:

- Any foregrounded process (via normal command or `fg`) that wants to read from the terminal (such as `cat` or `vim`) can.
- Any backgrounded process (via `command &` or `bg`) that tries to do so gets stopped.
- After a foreground process exits or is stopped, the shell is again able to read input from the terminal.
- Typing `ctrl-c` or `ctrl-z` during a foregrounded process's execution affects only that process and its children. (Your shell should *not* need to catch the signal and "forward it along", nor should a backgrounded process ever see the signal.)

7 Checking Your Work

Running your shell. The best way to check your work is to run your shell from the command line. Your initial testing should be done manually from the command line. Run your shell, type commands to it, and see if you can break it. Use it to run real programs!

Reference solution. The 64-bit Linux executable `tshref` is the reference solution for the shell. Run this program (on a 64-bit machine) to resolve any questions you have about how your shell should behave. Your shell should emit output that is identical to the reference solution (except for PIDs, of course, which change from run to run).

Once you are confident that your shell is working, then you can begin to use some tools that we have provided to help you check your work more thoroughly. (These are the same tools that the autograder will use when you submit your work for credit.)

Trace interpreter. We have provided a set of trace files (`trace*.txt`) that validate the correctness of your shell (the appendix section at the end of this handout describes each trace file briefly). Each trace file tests one shell feature. For example, does your shell recognize a particular built-in command? Does it respond correctly to the user typing a `ctrl-c`?

The `runtrace` program (the trace interpreter) interprets a set of shell commands specified by a single trace file:

```
linux> ./runtrace -h
Usage: runtrace -f <file> -s <shellprog> [-hV]
Options:
  -h                Print this message
  -s <shell>        Shell program to test (default ./tsh)
  -f <file>         Trace file
  -V                Be more verbose
```

The neat thing about the trace files is that they generate the same output you would have gotten had you run your shell interactively (except for an initial comment that identifies the trace). For example:

```
linux> ./runtrace -f trace05.txt -s ./tsh
#
# trace05.txt - Run a background job.
#
tsh> ./myspin1 &
[1] (15849) ./myspin1 &
tsh> quit
```

The lower-numbered trace files do very simple tests, and the higher-numbered tests do increasingly more complicated tests.

Shell driver. After you have used `runtrace` to test your shell on each trace file individually, then you are ready to test your shell with the shell driver. The `sdriver` program uses `runtrace` to run your shell on each trace file, compares the output to the output produced by the reference shell, displays the `diff` if they differ, and optionally sends the results to the Autolab server:

```
linux> ./sdriver -h
Usage: sdriver [-hV] [-s <shell> -t <tracenum> -i <iters>]
Options
  -h                Print this message.
  -i <iters>        Run each trace <iters> times (default 4)
  -s <shell>        Name of test shell (default ./tsh)
```

-t <n> Run trace <n> only (default all)
-V Be more verbose.

Running the driver without any arguments tests your shell on all of the trace files. To help detect race conditions in your code, the driver runs each trace multiple times. You will need to pass each of the tests to get credit for a particular trace:

```
linux> ./sdriver
Running 4 iters of trace00.txt
1. Running trace00.txt...
2. Running trace00.txt...
3. Running trace00.txt...
4. Running trace00.txt...
Running 4 iters of trace01.txt
1. Running trace01.txt...
2. Running trace01.txt...
3. Running trace01.txt...
4. Running trace01.txt...
Running 4 iters of trace02.txt
1. Running trace02.txt...
2. Running trace02.txt...
3. Running trace02.txt...
4. Running trace02.txt...

...

Running 4 iters of trace23.txt
1. Running trace23.txt...
2. Running trace23.txt...
3. Running trace23.txt...
4. Running trace23.txt...
Running 4 iters of trace24.txt
1. Running trace24.txt...
2. Running trace24.txt...
3. Running trace24.txt...
4. Running trace24.txt...
```

Summary: 25/25 correct traces

Use the optional -i argument to control the number of times the driver runs each trace file:

```
linux> ./sdriver -i 1
Running trace00.txt...
Running trace01.txt...
Running trace02.txt...
Running trace03.txt...

...

Running trace23.txt...
```

Running trace24.txt...

Summary: 25/25 correct traces

Use the optional `-t` argument to test a single trace file:

```
linux> ./sdriver -t 06
```

Running trace06.txt...

Success: The test and reference outputs for trace06.txt matched!

Note: The driver program runs the reference shell, which is a 64-bit binary, and thus will not run on a 32-bit machine.

8 Hints

- Read and understand every word of Chapter 8 (Exceptional Control Flow) and Chapter 11 (System-level I/O) in your textbook.
- Read the code in `tsh.c` carefully before you start. Understand the high-level control flow, get familiar with the defined global variables and the helper routines.
- Play with your shell by typing commands to it directly. Don't make the mistake of running the trace generator and driver immediately. Develop some familiarity and intuition about how your shell works before testing it with the automated tools.
- Only after you have tested your shell directly from the command and are fairly confident that it is correct should you start testing with the `runtrace` and driver programs.
- Use the trace files to guide the development of your shell. Starting with `trace00.txt`, make sure that your shell produces the *identical* output as the reference shell. Then move on to trace file `trace01.txt`, and so on.
- The `waitpid`, `kill`, `fork`, `execve`, `setpgid`, `tcsetpgrp`, `isatty`, `sigprocmask`, and `sigsuspend` functions will come in very handy. The `WUNTRACED` and `WNOHANG` options to `waitpid` will also be useful. Use `man` to check out the details about each function.
- One of the tricky parts of the assignment is deciding on the allocation of work between the `eval` and `sigchld_handler` functions when the shell is waiting for a foreground job to finish. For example, you will need to think about the various ways you can make your shell wait for a foreground process to exit (`waitpid` and `sigsuspend`) and where they should each be used.
- Remember that using `sleep` and/or busy-waiting are *never* acceptable ways of solving concurrency problems.
- Be careful about race conditions on the job list. Remember that you cannot make any assumptions about the order of execution of the parent and child after forking. In particular, you cannot assume that the child will still be running when the parent returns from the `fork`. In fact, our driver has code

that purposely introduces non-determinism in the order that the parent and child execute after forking. Also, remember that signal handlers run concurrently with the program and can interrupt it anywhere, unless you explicitly block the receipt of the signals.

- In `eval`, the parent must use `sigprocmask` to block `SIGCHLD` signals before it forks the child, and then unblock it, again using `sigprocmask` after it adds the child to the job list by calling `addjob`. Since children inherit the blocked vectors of their parents, the child must be sure to then unblock `SIGCHLD` before it execs the new program. The child should also restore the default handlers for the signals that are ignored by the shell.

The parent needs to block signals in this way in order to avoid race conditions (e.g., the child is reaped by `sigchld_handler` (and thus removed from the job list) *before* the parent calls `addjob`). Section 8.5.6 has details about the race conditions and how to block signals explicitly.

Please note that *blocking* signals is different from *ignoring* signals - when a signal is blocked, it will be received when it is unblocked, and when a signal is ignored, it will not be received. Your shell should ignore `SIGINT` and `SIGTSTP` but not `SIGCHLD`, and sometimes block `SIGCHLD` but not necessarily `SIGINT` or `SIGTSTP`.

9 Evaluation

Your score will be computed out of a maximum of 139 points based on the following distribution:

100 Correctness: 25 trace files at 4 pts each

24 Correctness: 6 secret trace files at 4 pts each

5 Correctness: Hand-graded use of `tcsetpgrp`

10 Style points. We expect you to have good comments and to check the return value of EVERY system call. We also expect you to break up large functions such as `eval` into smaller helper functions, to enhance readability and avoid duplicating code. Some advice about commenting:

- Do begin each routine with a block comment describing its role at a high level.
- Do preface related lines of code with a block comment.
- Do use consistent indenting and brace style.
- Do keep your comments and code within 80 character lines.
- Don't simply comment each line.

You should also follow other guidelines of good style, such as using a consistent indentation scheme, using descriptive variable names, and grouping logically related blocks of code with whitespace.

Please note that the driver program is not a perfect test for race conditions, and you should study your code very carefully to ensure that you have no race conditions. The TAs will read over the code and verify correctness, and reserve the right to deduct points after the end of the lab if that is necessary.

Your solution shell will be tested for correctness on a 64-bit fish machine (the Autolab server), using the same driver and trace files that were included in your handout directory. Your shell should produce **identical** output on these traces as the reference shell, with only two exceptions:

- The PIDs can (and will) be different.
- The output of the `/bin/ps` commands in `trace19.txt`, `trace20.txt`, and `trace21.txt` will be different from run to run. However, the running states of any `mysplit` processes in the output of the `/bin/ps` command should be identical.

The driver deals with all of these subtleties when it checks for correctness.

10 Hand In Instructions

- Make sure you have included your name and Andrew ID in the header comment of `tsh.c`.
- Hand in your `tsh.c` file for credit by uploading it to Autolab. You may hand in as often as you like. You will be graded on the **last** version you hand in.
- After you hand in, it takes a minute or two for the driver to run through multiple iterations of each trace file.
- We'll be using a sophisticated cheat checker that compares handins from this year and previous years. Please don't copy another student's code. Start early, and if you get stuck, come see your instructors for help.

Good luck!

Appendix: Trace Files

The trace driver runs an instance of your shell in a child process and communicates with the shell interactively in a way that mimics the behavior of a user. To test the behavior of your shell, the trace driver reads in trace files that specify shell line commands (that are actually sent to the shell) as well as a few special synchronization commands (that are interpreted by the driver when handling the shell process). The trace files may also reference a number of shell test programs to perform various functions, and you may refer to the code and comments of these test programs for more information.

The format of the trace files is as follows:

- The comment character is #. Everything to the right of it on a line is ignored.
- Each trace file is written so that the output from the shell shows exactly what the user typed. We do this by using the `/bin/echo` program, which not only tests the shell's ability to run programs, but also shows what the user typed. For example:

```
/bin/echo -e tsh\076 ./myspin1 \046
```

Note: octal `\076` is `>` and octal `\046` is `&`. These are special shell metacharacters that need to be escaped. This line represents `tsh> ./myspin1 &`, that is, a user trying to run `./myspin1` in the background.

- There are also a few special commands for synchronization between the job (your shell) and the parent process (the driver) and to send Unix signals from the parent to the job.

WAIT	Wait for a sync signal from the job over its synchronizing domain socket.
SIGNAL	Send a sync signal to the job over its synchronizing domain socket.
NEXT	Read and print all responses from the shell until you see the next shell prompt. This command is essential for synchronizing with the shell and mimics the way people wait until they see the shell prompt until they type the next string.

The following table describes what each trace file tests on your shell against the reference solution.

NOTE: this table is provided so that you can quickly get a high level picture about the testing traces. The explanation here is over-simplified. To understand what exactly each trace file does, you need to read the trace files.

trace00.txt	Properly terminate on EOF.
trace01.txt	Process built-in quit command.
trace02.txt	Run a foreground job that prints an environment variable.
trace03.txt	Run a synchronizing foreground job without any arguments.
trace04.txt	Run a foreground job with arguments.
trace05.txt	Run a background job.
trace06.txt	Run a foreground job and a background job.
trace07.txt	Use the jobs built-in command.
trace08.txt	Send fatal SIGINT to foreground job.
trace09.txt	Send SIGTSTP to foreground job.
trace10.txt	Send fatal SIGTERM (15) to a background job.
trace11.txt	Child sends SIGINT to itself.
trace12.txt	Child sends SIGTSTP to itself.
trace13.txt	Correct handling of SIGINT with foreground job present.
trace14.txt	Correct handling of SIGTSTP with foreground job present.
trace15.txt	Process bg built-in command (one job).
trace16.txt	Process bg built-in command (two jobs).
trace17.txt	Process fg built-in command (one job).
trace18.txt	Process fg built-in command (two jobs).
trace19.txt	Foreground child sends SIGINT to its own process group.
trace20.txt	Foreground child sends SIGTSTP to its own process group.
trace21.txt	Restart every stopped process in process group.
trace22.txt	I/O redirection (input).
trace23.txt	I/O redirection (input and output).
trace24.txt	I/O redirection (input and output, but different order).
secret-trace00.txt	Process built-in quit and jobs commands.
secret-trace01.txt	Send fatal SIGTERM to a background job.
secret-trace02.txt	Correct handing of SIGINT and SIGTSTP with foreground job.
secret-trace03.txt	Process bg and fg built-in commands.
secret-trace04.txt	Foreground child sends SIGINT and SIGTSTP to its own process group.
secret-trace05.txt	I/O redirection.